# Tensor networks in quantum information

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# What is it?

• A diagrammatic method for representing tensors

- Why is this useful?
- Gives a convenient way of bookkeeping how much entanglement there is between different parts of a quantum system
- Has resulted in many useful numerical and analytic methods for analysing interesting many-body quantum states.

#### Goals

1. Get comfortable with the notation

2. See how tensor network methods are used in quantum information/condensed matter

## Reference

- These slides use figures from the lecture notes by Bridgeman and Chubb:
  - Hand-waving and Interpretive Dance: An Introductory Course on Tensor Networks, J. Phys. A: Math. Theor. 50 223001 (2017) arXiv:1603.03039

• Much of what we will cover can also be found in lectures 1-3 in those notes.

#### Tensors



- Represent tensors as shapes
- Indices represented as legs

#### Combining tensors: tensor product

•  $[A \otimes B]_{i_1, \dots, i_r; j_1, \dots, j_s} \coloneqq A_{i_1, \dots, i_r} \cdot B_{j_1, \dots, j_s}$ 



#### Contracting indices: trace

•  $\left[\operatorname{Tr}_{x,y}A\right]_{i_1,\dots,i_{x-1},i_{x+1},\dots,i_{y-1},i_{y+1},\dots,i_r} \coloneqq \sum_{\alpha=1}^{d_x} A_{i_1,\dots,i_{x-1},\alpha,i_{x+1},\dots,i_{y-1},\alpha,i_{y+1},\dots,i_r}$ 



• Cyclic property of the trace:

$$\operatorname{Tr}(AB) = \begin{array}{c} \hline \\ A \\ \hline \\ B \\ \hline \\ B \\ \hline \\ A \\ \hline \\ B \\ \hline \\ B \\ \hline \\ A \\ \hline \\ B \\ \hline \\ A \\ \hline \\ B \\ \hline \\ B \\ \hline \\ A \\ \hline \\ B \\ \hline \\ B \\ \hline \\ A \\ \hline \\ B \\ \hline \\$$

#### Contracting indices: generally



Conventional	Einstein	$\operatorname{TNN}$
$\langle ec{x}, ec{y}  angle$	$x_{lpha}y^{lpha}$	x - y
Mec v	$M^{lpha}_{\ eta}v^{eta}$	-M-v
AB	$A^{lpha}_{\ eta}B^{eta}_{\ \gamma}$	
$\operatorname{Tr}(X)$	$X^{lpha}_{\ lpha}$	$\widehat{X}$

#### Grouping and splitting indices

• 
$$\mathbb{C}^{a_1 \times \dots \times a_n} \cong \mathbb{C}^{b_1 \times \dots \times b_m} \iff \prod_{i=1}^n a_i = \prod_{i=1}^m b_i$$

• E.g.  $\mathbb{C}^4 \cong \mathbb{C}^{2 \times 2} \cong \mathbb{C}^2 \otimes (\mathbb{C}^2)^*$ 

$$\begin{pmatrix} ac \\ ad \\ bc \\ bd \end{pmatrix} \cong \begin{pmatrix} a \\ b \end{pmatrix} \otimes \begin{pmatrix} c \\ d \end{pmatrix} \cong \begin{pmatrix} ac & ad \\ bc & bd \end{pmatrix}$$
$$= \bigcirc = \bigcirc \cong \bigcirc = \bigcirc = \bigcirc = \bigcirc$$

Turning tensors into matrices



- $T_{i_1\dots i_n;j_1\dots j_m} = T_{I;J}$
- $I = i_1 + \dim(i_1) i_2 + \dim(i_1) \dim(i_2) i_3 \dots$
- $J = j_1 + \dim(j_1) j_2 + \dim(j_1) \dim(j_2) j_3 \dots$
- E.g., if each index of T has dimension 2
  - $|i_1\rangle \otimes |i_2\rangle \cong |I\rangle$ , where  $i_1$  and  $i_2 \in \{0, 1\}, I \in \{0, 1, 2, 3\}$
  - $\langle j_1 | \otimes \langle j_2 | \otimes \langle j_3 | \cong \langle J |$ , where  $j_1, j_2$ , and  $j_3 \in \{0, 1\}$ ,  $J \in \{0, 1, 2, 3, 4, 5, 6, 7\}$



- Singular value decomposition:
  - $T_{I;J} = \sum_{\alpha} U_{I;\alpha} S_{\alpha;\alpha} V_{J;\alpha}^*$ , where  $U^{\dagger}U = I, S > 0$  is diagonal,  $V^{\dagger}V = I$
- For square matrices: Polar decomposition:
  - T = WP, where  $W = UV^{\dagger}$  is unitary,  $P = VSV^{\dagger} > 0$
- For Hermitian matrices: Spectral decomposition
  - $T = UDU^{\dagger}$ , where U is unitary and D is diagonal & real

#### Isometries vs Unitaries

- Unitaries map  $\mathbb{C}^d \to \mathbb{C}^d$ 
  - Written as  $U = \sum_{i=1}^{d} |v_i\rangle \langle e_i|$ , for some  $\mathbb{C}^d$  orthonormal bases  $\{|e_i\rangle\}$  and  $\{|v_i\rangle\}$
  - Have  $U^{\dagger}U = I = UU^{\dagger}$  over  $\mathbb{C}^d$
- Isometries  $\mathbb{C}^k \to \mathbb{C}^d$  for k < d
  - Written as  $W = \sum_{i=1}^{k} |v_i\rangle \langle e_i|$ , for some  $\mathbb{C}^k$ orthonormal basis  $\{|e_i\rangle\}$ , and  $\{|v_i\rangle\}$  a strict subset of orthonormal basis elements of  $\mathbb{C}^d$
  - Have  $W^{\dagger}W = I$  over  $\mathbb{C}^k$
  - Have  $WW^{\dagger} = \Pi_k$ , a rank k projector over  $\mathbb{C}^d$

#### Rank-k approximations to tensors

- 1. Treat the tensor as a matrix by grouping the legs into two sets (rows and columns)
- 2. Perform singular value decomposition
- 3. Set all but the largest k singular values equal to zero (known as k-trimming)
- Eckart-Young theorem:
  - Let  $X = USV^{\dagger}$ . Define  $X^{(k)} = US^{(k)}V^{\dagger}$  to be the *k*-trimmed version of *S*. Then  $||X X^{(k)}||_F \le ||X Y||_F$  for all rank-*k* matrices *Y*.
  - Note:  $||A||_F = \sqrt{\operatorname{Tr}(A^{\mathrm{T}}A)}$

## Contracting tensor networks



- Contracting general tensor networks is challenging
- Fortunately, we can reduce this to sequential matrix multiplication

# Bubbling



Important note: horizontal-first bubbling on a  $2 \times l$  lattice will result in a  $\chi^{l-1} \times \chi$  intermediate matrix

#### Bubbling: order matters!



Important note: vertical-first bubbling on a  $2 \times l$  lattice will result in no intermediate matrices larger than  $\chi^2 \times \chi$  or  $\chi \times \chi^2$ 

# 1D vs 2D networks





- There is no choice of bubbling that can avoid exponentially growing matrices on the 2D square lattice network.
- Contracting such networks is very hard.
  - Contains #P-complete problems
  - Approximating such contractions is Post-BQP hard

# TNN and quantum information

General quantum states

• 
$$|\psi\rangle = \sum_{i_1 i_2 \dots = 0}^{1} c_{i_1 i_2 \dots} |i_1 i_2 \dots \rangle$$

$$= \begin{array}{c} C_{\vec{l}} \\ \hline \end{array}$$

Product states

• 
$$|\psi\rangle = \bigotimes_{i=1}^{n} |\phi_i\rangle$$

$$\begin{array}{c|c} \phi_1 & \phi_2 & \phi_3 & \phi_4 \\ \hline & & & & & \\ \end{array}$$

# TNN and quantum information

• The Bell basis:

• 
$$\left| \Phi^{\pm} \right\rangle = \frac{\left| 00 \right\rangle \pm \left| 11 \right\rangle}{\sqrt{2}}, \left| \Psi^{\pm} \right\rangle = \frac{\left| 01 \right\rangle \pm \left| 10 \right\rangle}{\sqrt{2}}$$

• Refer to 
$$\left| \Phi^+ \right\rangle$$
 as  $\left| \Omega \right\rangle$ 

$$|\Omega\rangle = \frac{1}{\sqrt{2}} \begin{pmatrix} 1\\ 0\\ \cdot \\ 0\\ 1 \end{pmatrix}$$

Vectorise Matricise

$$\frac{1}{\sqrt{2}} \begin{pmatrix} 1 & 0\\ 0 & 1 \end{pmatrix} = I/\sqrt{2}.$$

- $|\Omega(I)\rangle = |\Phi^+\rangle$
- $|\Omega(X)\rangle = |\Psi^+\rangle$
- $|\Omega(Y)\rangle \propto |\Psi^-\rangle$
- $|\Omega(Z)\rangle = |\Phi^-\rangle$

$$\square \Omega(O) = \frac{1}{\sqrt{2}} \square O$$

 $\boxed{\Omega} = \frac{1}{\sqrt{2}}$ 

## Teleportation

• If Alice and Bob share an entangled resource state, and are allowed to perform classical communication, then Alice can perfectly transmit an arbitrary quantum state to Bob



## Purification

• Given a mixed state  $\rho$  over some Hilbert space  $\mathcal{H}_1$ , find a pure state  $|\psi\rangle$  from an extended Hilbert space  $\mathcal{H}_1 \otimes \mathcal{H}_2$  s.t.  $\rho = \mathrm{Tr}_2(|\psi\rangle\langle\psi|)$ 



# Stinespring dilation

- A quantum channel *E* that takes valid density matrices to valid density matrices is a completely-positive trace-preserving (CPTP) map.
- In the Kraus operator decomposition:

• 
$$\mathcal{E}(\rho) = \sum_{i} K_{i}^{\dagger} \rho K_{i}$$
, where  $\sum_{i} K_{i} K_{i}^{\dagger} = I$ 

In tensor network notation



Note the unfortunate dagger convention here

# Stinespring dilation

- The Stinespring dilation theorem says that any CPTP map  ${\mathcal E}$  can be made by
- 1. Embedding the system in a larger Hilbert space
- 2. Evolving the enlarged system under unitary dynamics
- 3. Reducing back to the original system via partial trace
- To see this, we use that the tensor  $K_i^{\dagger}$  is an isometry



## Stinespring dilation







 $= \operatorname{Tr}_{1}\left[ V^{\dagger}\left( \rho \otimes |0\rangle \langle 0| \right) V \right]$ 

# The Schmidt decomposition

- Given a bipartite quantum state
  - $|\psi\rangle_{AB} = \sum_{i=1}^{d} \sum_{j=1}^{d} c_{ij} |i\rangle \otimes |j\rangle$
- We can always write it as
  - $|\psi\rangle_{AB} = \sum_{i=1}^{\chi \le d} s_i |u(i)\rangle \otimes |v(j)\rangle$
  - $\chi$  is known as the Schmidt rank
  - Entanglement entropy  $S = -\sum_{i=1}^{\chi} s_i \log s_i \le \log \chi$

$$\begin{array}{c|c} C_{ij} & \xrightarrow{\text{SVD}} & u & -s & -v^{\dagger} \\ \hline & & & & & & & \\ \hline & & & & & & & \\ \end{array}$$

# The geometry of entanglement

- Under special circumstances, a many-body quantum system will have structured entanglement entropy. We can see this using tensor network notation.
- Example 1: Product state  $|\psi\rangle = \bigotimes_{i=1}^{n} |\phi_i\rangle$

$$\phi_1 \phi_2 \phi_3 \phi_4$$

• Example 2: Entangled pairs  $|\psi\rangle = \bigotimes_{i=1}^{n} |\varphi_i\rangle_{AB}$ 



. .

# The geometry of entanglement

• Example 3: *d*-dimensional regular lattices





Area law for entanglement  $S \sim \log \chi^{d-1}$ 

# The geometry of entanglement

• Example 4: other lattices

Extra layer gives longer range interactions



Area law for entanglement  $S \sim \log L(A)$ 



## Connection to physics

- Choose a state out of the many-body Hilbert space. Generically, it will have maximal entanglement entropy with respect to any bipartition.
- States that are considered physically reasonable correspond to low-energy sectors of Hamiltonians that are local with respect to a *d*-dimensional lattice and have *unique* ground states (or possibly bounded degeneracy).

## Area law conjecture

- Given the unique ground state of a *gapped*, local Hamiltonian on a *d*-dimensional lattice, that state satisfies a boundary law with respect to entanglement entropy.
- Proven in 1D
- Consequences in Tensor networks: We can find a good approximation to the ground state using tensors that look like:

#### Matrix product state approximations

