

Ph 452 Lecture 10  
Quantum Error Correction

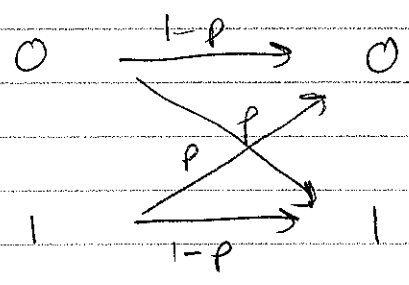
more errors for qubits  
Need to understand

Noisy Channel

Hello? Hello? Hello?

Redundancy helps!

Binary symmetric channel:



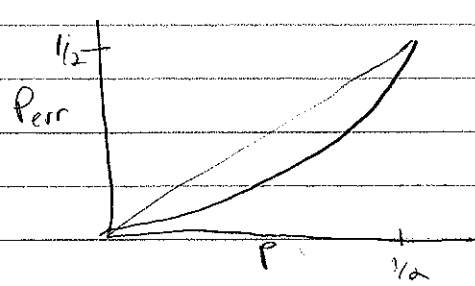
$$P = \begin{pmatrix} 1-p & p \\ p & 1-p \end{pmatrix} \quad \text{Pointe}$$

p is unknown!  
(otherwise could apply  $P^{-1}$ )

3-bit repetition code

0 → 000  
1 → 111

$$\begin{aligned} \text{Majority vote} \rightarrow P_{err} &= \binom{3}{0} p^3 + \binom{3}{2} p^2(1-p) \\ &= p^3 + 3p^2 - 3p^3 \\ &= 3p^2 - 2p^3 \end{aligned}$$



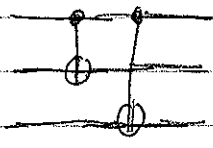
Improvement if  $p < 1/2$ !

Ph 452 Lecture 10

Reversible encoding:

$$000 \rightarrow 000$$

$$100 \rightarrow 111$$



What about pbits?

Encoding:

$$\vec{p} \otimes \vec{0} \otimes \vec{0} \mapsto \vec{p} \otimes \vec{p} \otimes \vec{p}$$

Impossible!

No-cloning Thm: There is no pgate that acts as  $\vec{p} \otimes \vec{0} \mapsto \vec{p} \otimes \vec{p}$  for all  $\vec{p} := (p_0, p_1)$

Proof: Spse there was such a P.

$$\text{Then } P(\vec{0} \otimes \vec{0}) = \vec{0} \otimes \vec{0}$$

$$P(\vec{1} \otimes \vec{1}) = \vec{1} \otimes \vec{1}$$

$$\text{By linearity, } P\left(\begin{pmatrix} 1/2 \\ 1/2 \end{pmatrix} \otimes \vec{0}\right) = \frac{1}{2} P(\vec{0} \otimes \vec{0}) + \frac{1}{2} P(\vec{1} \otimes \vec{0})$$

$$= \frac{1}{2} (\vec{0} \otimes \vec{0} + \vec{1} \otimes \vec{1})$$

$$= \begin{pmatrix} 1/2 \\ 0 \\ 0 \\ 1/2 \end{pmatrix}$$

$$\text{But } \begin{pmatrix} 1/2 \\ 1/2 \end{pmatrix} \otimes \begin{pmatrix} 1/2 \\ 1/2 \end{pmatrix} = \frac{1}{4} \begin{pmatrix} 1 \\ 1 \\ 1 \\ 1 \end{pmatrix} ! \quad \square$$

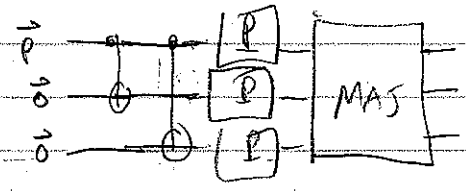
→ Same for (qn) gates: No U can map  $|\psi\rangle|0\rangle \mapsto |\psi\rangle|\psi\rangle \forall |\psi\rangle$ .

Ph 452 Lecture 10

Solution: Encode  $p$  bits with same circuit used for bits!

$$\vec{p} = \begin{pmatrix} p_0 \\ p_1 \end{pmatrix} \begin{array}{c} \text{---} \\ \oplus \\ \text{---} \\ \oplus \\ \text{---} \end{array} \begin{pmatrix} p_0 \\ 0 \\ \vdots \\ 0 \\ p_1 \end{pmatrix} = p_0 \vec{0} \otimes \vec{0} \otimes \vec{0} + p_1 \vec{1} \otimes \vec{1} \otimes \vec{1}$$

Full encode-error-correct circuit



With a slight abuse of notation, I'll use direct notation

$$MAJ = |000\rangle\langle 000| + |000\rangle\langle 001| + |000\rangle\langle 010| + |000\rangle\langle 100| + |111\rangle\langle 111| + |111\rangle\langle 110| + |111\rangle\langle 101| + |111\rangle\langle 011|$$

Not a  $p$  gate: both 001, 000 map to 000, ergo  
 Not a classical measurement: MAJ is not diagonal.

MAJ is a combination of  $p$  gates + measurements

Method 1: Measure 3 bits, then correct

$$\begin{array}{ll} \Pi_0 = |000\rangle\langle 000| & \Pi_3 = |011\rangle\langle 011| \\ \Pi_1 = |001\rangle\langle 001| & \vdots \\ \Pi_2 = |010\rangle\langle 010| & \Pi_7 = |111\rangle\langle 111| \end{array}$$

If outcome 0, apply  $P_0 = I \otimes I \otimes I$   
 " " 1, "  $P_1 = I \otimes I \otimes X$   
 " " 2, "  $I \otimes X \otimes I$   
 " " 3, "  $X \otimes I \otimes I$   
 " "  $\vdots$   
 " " 7,  $P_7 = I \otimes I \otimes I$

$$\vec{p}' \mapsto P_k \Pi_k \vec{p}'$$

w/ prob  $\sum_i (\Pi_i \vec{p})_i$

Ph 452 Lecture 10

Method 1 fails on pbits:

E.g., suppose  $p=0$  so  $P=I$ , Then

$$\vec{p} := p_0 |000\rangle + p_1 |111\rangle \mapsto \begin{cases} |000\rangle & \text{probability } p_0 \\ |111\rangle & \text{" } p_1 \end{cases}$$

But neither  $|000\rangle$  nor  $|111\rangle$  are  $\vec{p}$ !

"Measurement collapses  $\vec{p}$  into one of two states, neither of which is  $\vec{p}$ "

Method 2: Measure the error, not the data

e.g., if in state  $\begin{matrix} \text{parity} \\ \boxed{010} \\ \text{parity} \end{matrix}$  - 2 parities determine error!

$$\Pi_0 = |000\rangle\langle 000| + |001\rangle\langle 001| + |110\rangle\langle 110| + |111\rangle\langle 111|$$

$$\Pi_1 = I - \Pi_0$$

$$\Pi'_0 = |000\rangle\langle 000| + |100\rangle\langle 100| + |011\rangle\langle 011| + |111\rangle\langle 111|$$

$$\Pi'_1 = I - \Pi'_0$$

If outcome 00, apply  $P_0 = I \otimes I \otimes I$  }  $\vec{p}' \mapsto P_{ij} \Pi_i \Pi_j \vec{p}'$   
 " " 01, "  $P_1 = I \otimes I \otimes X$  }  
 " " 10, "  $P_2 = X \otimes I \otimes I$  } w/prob.  $\sum_k (P_k \Pi_k \vec{p}')_k$   
 " " "  $P_3 = I \otimes X \otimes I$  }

Ph 452 Lecture 10

Method 2 works on bits!

$$\begin{aligned}
 \vec{p} \otimes \vec{0} \otimes \vec{0} &= p_0 \vec{0} \otimes \vec{0} \otimes \vec{0} + p_1 \vec{1} \otimes \vec{1} \otimes \vec{1} \\
 &\xrightarrow{P^{\otimes 3}} \\
 & p_0 (1-p)^3 |000\rangle \\
 & + p_0 (1-p)^2 p (|100\rangle + |010\rangle + |100\rangle) \\
 & + p_0 (1-p) p^2 (|011\rangle + |101\rangle + |110\rangle) \\
 & + p_0 p^3 |111\rangle \\
 & + p_1 (1-p)^3 |111\rangle \\
 & + p_1 (1-p)^2 p (|110\rangle + |101\rangle + |101\rangle) \\
 & + p_1 (1-p) p^2 (|100\rangle + |1010\rangle + |1001\rangle) \\
 & + p_1 p^3 |1000\rangle
 \end{aligned}$$

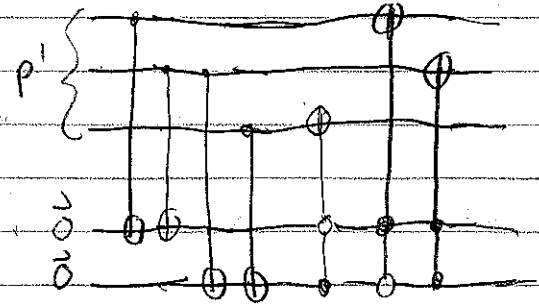
To do this coherently, need to entangle bits ahead of the syndrome.

$$\begin{aligned}
 &\xrightarrow{R_0, \Pi_1, \Pi_2} \\
 & \left\{ \begin{aligned} &(p_0 |1000\rangle + p_1 |1111\rangle) \quad \text{prob } (1-p)^3 + 3(1-p)^2 p = 1 - p_{err} \\ &(p_1 |1000\rangle + p_0 |1111\rangle) \quad \text{prob } 3(1-p) p^2 + p^3 = p_{err} \end{aligned} \right.
 \end{aligned}$$

vs. no encoding:

$$p_0 |0\rangle + p_1 |1\rangle \xrightarrow{P} (p_0 |0\rangle + p_1 |1\rangle)(1-p) + (p_1 |0\rangle + p_0 |1\rangle)p$$

Can also do method 2 w/o measurement:



$$\begin{aligned}
 &(p_0 |1000\rangle + p_1 |1111\rangle)(1 - p_{err}) \\
 &+ (p_1 |1000\rangle + p_0 |1111\rangle) p_{err}
 \end{aligned}$$

control on "0", not "1"

Ph 452 Lecture 10

This is as good a time as any to introduce...

Observables

Measurements of physical quantities have units

E.g.  $\Pi_0 = |0\rangle\langle 0|$ ,  $\Pi_1 = |1\rangle\langle 1|$   
 $f(0) = +\hbar/2$       $f(1) = -\hbar/2$

measures electron spin (units = angular momentum)

$A = \sum_i f(i) \Pi_i$  is an observable associated with  $\{\Pi_i\}$  taking distinct values  $\{f(i)\}$ .

Abstractly, any  $\{f(i)\}$  defines an observable, even w/o units:

E.g.  $f(0) = +1$ ,  $f(1) = -1$

$A = \sum (+1)|0\rangle\langle 0| + (-1)|1\rangle\langle 1| = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} = Z$  !

Similarly,  $X = \sum |+\rangle\langle +| - |-\rangle\langle -|$   
 $Y = \sum |+\rangle\langle +| - |-\rangle\langle -|$  ] Not classical observables because projectors involved are not diagonal matrices.

Expected value:

$\langle A \rangle = \sum p_i f(i)$   
 $= \sum \langle \psi | \Pi_i | \psi \rangle f(i)$   
 $= \langle \psi | \sum \Pi_i f(i) | \psi \rangle$   
 $= \langle \psi | A | \psi \rangle$

Ph 452 Lecture 10

Example: Parity observable for 2 qubits

$i, j \in \{1, 2\}$   
 $\#(I-Z)$

$$\Pi_0 = |00\rangle\langle 00| + |11\rangle\langle 11|$$

$$\Pi_1 = |01\rangle\langle 01| + |10\rangle\langle 10|$$

$$Z \otimes Z = (|0\rangle\langle 0| - |1\rangle\langle 1|) \otimes (|0\rangle\langle 0| - |1\rangle\langle 1|)$$

$$= (|00\rangle\langle 00| + |11\rangle\langle 11|) - (|01\rangle\langle 01| + |10\rangle\langle 10|)$$

Got this far  $\rightarrow$  An observable for parity!

Back to 3-bit repetition code

MAJ: ① Measure  $Z \otimes Z \otimes I$ ,  $(\pm 1)$   
 $I \otimes Z \otimes Z$ ,  $(\pm 1)$  ← "syndrome for the error"

Shortcut:  $ZZI, IZZ$  (Don't confuse with matrix multi.)

② Apply correction ("recovery")  $XII, IXI, IIX$  as appropriate.

$ZZI, IZZ$  & their products are called the stabilizer of the code

$$ZZI (p_0 |000\rangle + p_1 |111\rangle) = (p_0 |000\rangle + p_1 |111\rangle)$$

$$IZZ ( \quad ) = ( \quad )$$

$$ZIZ ( \quad ) = ( \quad )$$

$$III ( \quad ) = ( \quad )$$

Next time: Encoding & correcting errors in qubits